

Sorting Algorithms

The 3 sorting methods discussed here all have wild signatures. For example,

```
public static <E extends Comparable<? super E>>void BubbleSort(E[] array )
```

The underlined portion is a *type bound*. This says that the generic type E used as the base type of the array must implement or extend a superclass that implements, the Comparable interface (which says that E has a compareTo(E) method. See the discussion in Weiss of wildcards and type bounds,p. 151-154.

In less generic examples you probably don't need this. If you are writing BubbleSort to sort strings its signature could just be

```
public static void BubbleSort(String[] array)
```

BubbleSort

BubbleSort makes repeated passes through the array, interchanging successive elements that are out of order. When no changes are made in a pass the array is sorted.

```
public static <E extends Comparable<? super E>>void BubbleSort(E[]  
array ) {  
    boolean sorted = false;  
    int highest = array.length-1;  
    while (!sorted) {  
        sorted = true;  
        for (int i = 0; i < highest; i++) {  
            if (array[i].compareTo(array[i+1]) > 0) {  
                E buffer = array[i];  
                array[i] = array[i+1];  
                array[i+1] = buffer;  
                sorted = false;  
            }  
        }  
        highest -= 1;  
    }  
}
```

Original data

33	12	45	17	23	52	24
12	33	17	23	45	24	52
12	17	23	33	24	45	52
12	17	23	24	33	45	52
12	17	23	24	33	45	52

Each row shows the result of a pass through the previous row, flipping consecutive elements that are out of order.

The first pass through the list does $(n-1)$ comparisons. That pass puts the largest element into its proper location at the last spot in the list, so the next pass does $(n-2)$ comparisons. Altogether we do at most

$$(n-1)+(n-2)+\dots+1 = n(n-1)/2$$

comparisons. For each comparison we do at most 1 interchange, which takes 3 assignment statements. This means BubbleSort is worst-case $O(n^2)$.

Note that the best case for BubbleSort is when the data is already sorted; only one pass is then needed and the running time is $O(n)$. Of course, if you knew the data was already sorted there wouldn't be a lot of point in calling BubbleSort ...

SelectionSort

SelectionSort finds the smallest element and puts it at position 0, the smallest remaining element and puts it at position 1, etc.


```
public static <E extends Comparable<? super E>>void
    SelectionSort(E[] array ) {
    for (int i=0; i < array.length-1; i++ ) {
        // find the index of the smallest remaining element
        int small = i;
        for (int j = i+1; j < array.length; j++) {
            if (array[j].compareTo(array[small]) < 0)
                small = j;
        }
        // put the smallest remaining element at position i
        E buffer = array[i];
        array[i] = array[small];
        array[small]= buffer;
    }
}
```

Original data

33	12	45	17	23	52	24
----	----	----	----	----	----	----

12	33	45	17	23	52	24
----	----	----	----	----	----	----

12	17	45	33	23	52	24
----	----	----	----	----	----	----

12	17	23	33	45	52	24
----	----	----	----	----	----	----

12	17	23	24	45	52	33
----	----	----	----	----	----	----

12	17	23	24	33	52	45
----	----	----	----	----	----	----

12	17	23	24	33	45	52
----	----	----	----	----	----	----

The element put in its final location is in blue.

Selection sort does $(n-1)$ passes. The first one does $(n-1)$ comparisons; the second $(n-2)$ comparisons, and so forth. There are a total of

$$(n-1) + (n-2) + (n-3) + \dots + 1 = n(n-1)/2$$

comparisons. This is very similar to BubbleSort, only instead of interchanging elements of the array, which takes 3 assignments, here each comparison results in at most one integer assignment. Both are worst-case $O(n^2)$, but in specific examples SelectionSort usually runs somewhat faster.

Question: Suppose you use SelectionSort on an array of size n that is already sorted. How many comparisons will the sorting algorithm do?

- A. None
- B. 1
- C. $O(n)$
- D. $O(n^2)$

Answer D: $O(n^2)$.

Unlike BubbleSort, SelectionSort doesn't have a quick way out if the data is already sorted; it always does $n*(n-1)/2$ comparisons.

InsertionSort

InsertionSort maintains a sorted portion of the array (the front) and inserts elements from the unsorted portion into it.

```
public static <E extends Comparable<? super E>>void
    InsertionSort(E[] array ) {
    for (int p = 1; p < array.length; p++) {
        // p is the start of the unsorted portion
        E item = array[p];
        int j ;
        for ( j=p; j > 0 && item.compareTo(array[j-1]) < 0; j--)
            array[j] = array[j-1];
        array[j]= item;
    }
}
```

Original data

33	12	45	17	23	52	24
----	----	----	----	----	----	----

12	33	45	17	23	52	24
----	----	----	----	----	----	----

12	33	45	17	23	52	24
----	----	----	----	----	----	----

12	17	33	45	23	52	24
----	----	----	----	----	----	----

12	17	23	33	45	52	24
----	----	----	----	----	----	----

12	17	23	33	45	52	24
----	----	----	----	----	----	----

12	17	23	24	33	45	52
----	----	----	----	----	----	----

The sorted portion of the array is in blue.

It is easy to see that InsertionSort is no worse than $O(n^2)$ -- the outer loop runs n times, and the inner loop also takes at most n steps -- n steps done n times gives a total of n^2 steps.

The worst case is when the data is reverse-sorted (biggest to smallest); the first pass does 1 comparison, the second 2, and so forth.

Altogether this does $1+2+3+\dots+(n-1) = n(n-1)/2$ comparisons.

Question: Suppose you use InsertionSort on an array of size n that is already sorted. How many comparisons will the sorting algorithm do?

- A. None
- B. 1
- C. $O(n)$
- D. $O(n^2)$

Answer C: $O(n)$

If the data is already sorted, each pass does only one comparison and one assignment statement, so the algorithm runs in $O(n)$ steps.

InsertionSort is a good choice if you have a small amount of data to sort; it tends to be faster than the other simple sorts and is easy to implement.

If you want to sort data the size of the NY phone book, InsertionSort is a terrible choice. There are sorting algorithms that are $O(n \cdot \log(n))$, which is vastly better than $O(n^2)$ when n is large.